# On First-Order Knowledge Compilation 

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## Overview

1. Why first-order model counting?
2. Why first-order model counters?
3. What first-order circuit languages?
4. How first-order knowledge compilation?
5. Perspectives ...

# Why do we need <br> first-order model counting? 

## Uncertainty in AI

Probability Distribution =

Qualitative
+

Quantitative

## Probabilistic Graphical Models

## Probability Distribution

## =

Graph Structure
+

Parameterization

## Probabilistic Graphical Models

## Probability Distribution

## Graph Structure

+ 

Parameterization

$+$

| rain | $\operatorname{Pr}($ sun $\mid$ rain $)$ |
| :---: | :---: |
| T | 0.1 |
| F | 0.6 |


| rain | sun | $\operatorname{Pr}$ (rainbow $\mid$ rain, sun) |
| :---: | :---: | :---: |
| T | T | 0.9 |
| T | F | 0.05 |
| F | T | 0.05 |
| F | F | 0 |
| $\operatorname{Pr}($ rain $)$ |  |  |
| 0.2 |  |  |

## Weighted Model Counting

## Probability Distribution

SAT Formula

Weights

## Weighted Model Counting

## Probability Distribution

SAT Formula
+

Weights

Rain $\Rightarrow$ Cloudy
Sun $\wedge$ Rain $\Rightarrow$ Rainbow

$+$<br>$w($ Rain $)=1$<br>$w(\neg$ Rain $)=2$<br>$w($ Cloudy $)=3$<br>$w(\neg$ Cloudy $)=5$

## Beyond NP Pipeline for \#P


[Chavira 2006, Chavira 2008, Sang 2005, Fierens 2015]

## Generalized Perspective

Probability Distribution

## Logic

$+$
Weights

## Generalized Perspective

Probability Distribution

## Logic

$+$
Weights

## Logical Syntax Model-theoretic Semantics

$$
+
$$

Weight function w(.)
Factorized
$\operatorname{Pr}($ model $) \propto \Pi_{\mathrm{i}} \mathrm{w}\left(\mathrm{x}_{\mathrm{i}}\right)$

## First-Order Model Counting

## Probability Distribution

## = <br> First-Order Logic <br> $+$ <br> Weights



## First-Order Model Counting

## Probability Distribution

## First-Order Logic

Weights
Smokes $(\mathrm{x}) \wedge$ Friends $(\mathrm{x}, \mathrm{y})$ $\Rightarrow$ Smokes(y)

$$
+
$$

$w($ Smokes(a) $)=1$
$w(\neg$ Smokes(a)) $=2$
w( Smokes(b))=1
$w(\neg$ Smokes(b)) $=2$
$w($ Friends $(a, b))=3$ $\mathrm{w}(\neg$ Friends $(\mathrm{a}, \mathrm{b}))=5$

## Probabilistic Programming

## Probability Distribution

## =

Logic Programs

$$
+
$$

Weights

## Probabilistic Programming

## Probability Distribution

Logic Programs

$$
+
$$

Weights
path $(X, Y)$ :-
edge $(X, Y)$. $\operatorname{path}(X, Y)$ :edge $(X, Z)$, path $(Z, Y)$.
+

[Fierens 2015]

## Weighted Model Integration

## Probability Distribution

## = <br> SMT(LRA) <br> Weights

## Weighted Model Integration

## Probability Distribution

SMT(LRA)

Weights

$$
\begin{aligned}
& 0 \leq \text { height } \leq 200 \\
& 0 \leq \text { weight } \leq 200 \\
& 0 \leq \text { age } \leq 100 \\
& \text { age }<1 \Rightarrow \\
& \quad \text { height+weight } \leq 90
\end{aligned}
$$

$$
+
$$

$$
w(\text { height }) \text { ) }=\text { height }-10
$$

$$
\text { w( } \neg \text { height })=3^{*} \text { height² }^{2}
$$

$$
w(\neg \text { weight })=5
$$

## Beyond NP Pipeline for \#P/\#P ${ }_{1}$


[Van den Broeck 2011, 2013, Gogate 2011, Gribkoff 2014]

## First-Order Model Counting

Model $=$ solution to first-order logic formula $\Delta$

```
\Delta= \foralld (Rain(d)
    => Cloudy(d))
```

Days $=\{$ Monday $\}$

## First-Order Model Counting

Model = solution to first-order logic formula $\Delta$


FOMC = 3

## Weighted First-Order Model Counting

Model = solution to first-order logic formula $\Delta$

| $\Delta=\forall d$ |
| :--- |
| $($ Rain $(\mathrm{d})$ |
| $\Rightarrow \operatorname{Cloudy}(\mathrm{d}))$ |

Days $=\{$ Monday Tuesday\}

| Rain(M) | Cloudy(M) |
| :---: | :---: |
| T | T |
| T | F |
| F | T |
| F | F |


| Rain(T) | Cloudy(T) |
| :---: | :---: |
| T | T |
| T | T |
| T | T |
| T | T |


| Model? |
| :---: |
| Yes |
| No |
| Yes |
| Yes |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $T$ | $F$ |
| :---: | :---: |
| $T$ | $F$ |
| $T$ | $F$ |
| $T$ | $F$ |


| No |
| :---: |
| No |
| No |
| No |


| T | T |
| :---: | :---: |
| T | F |
| F | T |
| F | F |


| $F$ | $T$ |
| :---: | :---: |
| $F$ | $T$ |
| $F$ | $T$ |
| $F$ | $T$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $F$ |
| :---: | :---: |
| $F$ | $F$ |
| $F$ | $F$ |
| $F$ | $F$ |

Yes
No
Yes
Yes

## Weighted First-Order Model Counting

Model = solution to first-order logic formula $\Delta$

| $\Delta=\forall d$ |
| :--- |
| $($ Rain $(\mathrm{d})$ |
| $\Rightarrow \operatorname{Cloudy}(\mathrm{d}))$ |

$$
\begin{aligned}
\text { Days }= & \{\text { Monday } \\
& \text { Tuesday }\}
\end{aligned}
$$

| Rain(M) | Cloudy(M) |
| :---: | :---: |
| T | T |
| T | F |
| F | T |
| F | F |


| Rain(T) | Cloudy(T) | Model? |
| :---: | :---: | :---: | :---: |
| T | T | Yes |
| T | T | No |
| T | T | Yes |
| T | T | Yes |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| T | F |
| :---: | :---: |
| T | F |
| T | F |
| T | F |


| No |
| :---: |
| No |
| No |
| No |


| T | T |
| :---: | :---: |
| T | F |
| F | T |
| F | F |


| $F$ | $T$ |
| :---: | :---: |
| $F$ | $T$ |
| $F$ | $T$ |
| $F$ | $T$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $F$ |
| :---: | :---: |
| $F$ | $F$ |
| $F$ | $F$ |
| $F$ | $F$ |

Yes
No
Yes
Yes

## Weighted First-Order Model Counting

Model = solution to first-order logic formula $\Delta$

| $\Delta=\forall d$ |
| :--- |
| $($ Rain $(\mathrm{d})$ |
| $\Rightarrow \operatorname{Cloudy}(\mathrm{d}))$ |


| Days $=\{$ Monday | $F$ | $T$ |
| :--- | :--- | :--- |
|  | $F$ | $F$ |


| $\operatorname{Rain}(\mathrm{T})$ | $\mathbf{C l o u d y}(\mathbf{T})$ |
| :---: | :---: |
| T | T |
| T | T |
| T | T |
| T | T |


| Model? |
| :---: |
| Yes |
| No |
| Yes |
| Yes |


| Weight |  |
| :---: | ---: |
| $1 * 1 * 3 * 3=$ | 9 |
| $2 * 1 * 3 * 3=$ | 18 |
| $2 * 1 * 5 * 3=$ | 30 |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $T$ | $F$ |
| :---: | :---: |
| $T$ | $F$ |
| $T$ | $F$ |
| $T$ | $F$ |


| No |
| :---: |
| No |
| No |
| No |


| 0 |
| :--- |
| 0 |
| 0 |
| 0 |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $T$ |
| :---: | :---: |
| $F$ | $T$ |
| $F$ | $T$ |
| $F$ | $T$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |

$1 * 2 * 3 * 3=18$
0
$2 * 2 * 3 * 3=36$
$2 * 2 * 5 * 3=60$

| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $F$ |
| :---: | :---: |
| $F$ | $F$ |
| $F$ | $F$ |
| $F$ | $F$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |
| \#SAT =9 |

$1 * 2 * 3 * 5=30$
0
$2 * 2 * 3 * 5=60$
$2 * 2 * 5 * 5=100$

## Weighted First-Order Model Counting

Model = solution to first-order logic formula $\Delta$
$\Delta=\forall d$ (Rain(d)
$\Rightarrow \operatorname{Cloudy}(\mathrm{d}))$

| Days $=\{$ Monday | $F$ | $T$ |
| :--- | :--- | :--- |
|  | $F$ | $F$ |


| $\operatorname{Rain}(\mathrm{T})$ | $\mathbf{C l o u d y}(\mathbf{T})$ |
| :---: | :---: |
| T | T |
| T | T |
| T | T |
| T | T |


| Model? |
| :---: |
| Yes |
| No |
| Yes |
| Yes |


| Weight |
| :---: |
| $1 * 1^{*} 3 * 3=$ |
|  |
| $2 * 1 * 3 * 3=$ |
| $2 * 1 * 5 * 3=$ |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $T$ | $F$ |
| :---: | :---: |
| $T$ | $F$ |
| $T$ | $F$ |
| $T$ | $F$ |


| No |
| :---: |
| No |
| No |
| No |


|  |
| ---: |
|  |
|  |


| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $T$ |
| :---: | :---: |
| $F$ | $T$ |
| $F$ | $T$ |
| $F$ | $T$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |

$1 * 2 * 3 * 3=18$
0
$2 * 2 * 3 * 3=36$
$2 * 2 * 5 * 3=60$

| $T$ | $T$ |
| :---: | :---: |
| $T$ | $F$ |
| $F$ | $T$ |
| $F$ | $F$ |


| $F$ | $F$ |
| :---: | :---: |
| $F$ | $F$ |
| $F$ | $F$ |
| $F$ | $F$ |


| Yes |
| :---: |
| No |
| Yes |
| Yes |


| $1 * 2 * 3 * 5=30$ |
| ---: |
| 0 |
| $2 * 2 * 3 * 5=60$ |
| $2 * 2 * 5 * 5=100$ |
| + WFOMC $=\mathbf{3 6 1}$ |

# Why do we need first-order model counters? 

## A Simple Reasoning Problem



- 52 playing cards
- Let us ask some simple questions


## A Simple Reasoning Problem



Probability that Card1 is Hearts?

## A Simple Reasoning Problem



Probability that Card1 is Hearts?
1/4

## A Simple Reasoning Problem



Probability that Card1 is Hearts
given that Card1 is red?

## A Simple Reasoning Problem



Probability that Card1 is Hearts given that Card1 is red?

1/2

## A Simple Reasoning Problem



Probability that Card52 is Spades given that Card1 is QH?

## A Simple Reasoning Problem



Probability that Card52 is Spades given that Card1 is QH?

13/51

## A Simple Reasoning Problem



Probability that Card1 is Hearts?

## A Simple Reasoning Problem



Probability that Card1 is Hearts?
1/4

## A Simple Reasoning Problem



Probability that Card52 is Spades given that Card1 is QH?

## A Simple Reasoning Problem



Probability that Card52 is Spades given that Card1 is QH?

13/51


## Model distribution by FOMC:

$$
\begin{array}{r}
\forall \mathrm{p}, \exists \mathrm{c}, \operatorname{Card}(\mathrm{p}, \mathrm{c}) \\
\forall \mathrm{c}, \exists \mathrm{p}, \operatorname{Card}(\mathrm{p}, \mathrm{c}) \\
\forall \mathrm{p}, \forall \mathrm{c}, \forall \mathrm{c}^{\prime}, \operatorname{Card}(\mathrm{p}, \mathrm{c}) \wedge \operatorname{Card}\left(\mathrm{p}, \mathrm{c}^{\prime}\right) \Rightarrow \mathrm{c}=\mathrm{c}^{\prime}
\end{array}
$$

## Beyond NP Pipeline for \#P

Reduce to propositional model counting:

## Beyond NP Pipeline for \#P

Reduce to propositional model counting:

$$
\begin{aligned}
& \Delta=\operatorname{Card}\left(A \vee, p_{1}\right) \vee \ldots v \operatorname{Card}\left(2 \boldsymbol{\imath}, \mathrm{p}_{1}\right) \\
& \operatorname{Card}\left(A \vee, \mathrm{p}_{2}\right) \vee \ldots \vee \operatorname{Card}\left(2 \&, \mathrm{p}_{2}\right) \\
& \operatorname{Card}\left(\mathrm{A} \boldsymbol{\vee}, \mathrm{p}_{1}\right) \vee \ldots \vee \operatorname{Card}\left(\mathrm{A} \boldsymbol{\vee}, \mathrm{p}_{52}\right) \\
& \operatorname{Card}\left(K \vee, p_{1}\right) \vee \ldots v \operatorname{Card}\left(K \vee, p_{52}\right) \\
& \neg \operatorname{Card}\left(\mathrm{A} \vee, \mathrm{p}_{1}\right) \vee \neg \operatorname{Card}\left(\mathrm{A} \vee, \mathrm{p}_{2}\right) \\
& \neg \operatorname{Card}\left(A \vee, p_{1}\right) \vee \neg \operatorname{Card}\left(A \vee, p_{3}\right)
\end{aligned}
$$

## Beyond NP Pipeline for \#P

Reduce to propositional model counting:

$$
\begin{aligned}
& \Delta=\operatorname{Card}\left(\mathrm{A} \mathbf{v}, \mathrm{p}_{1}\right) \vee \ldots \vee \operatorname{Card}\left(2 \boldsymbol{*}, \mathrm{p}_{1}\right) \\
& \operatorname{Card}\left(A \vee, p_{2}\right) \vee \ldots v \operatorname{Card}\left(2 \&, \mathrm{p}_{2}\right) \\
& \operatorname{Card}\left(\mathrm{A} \boldsymbol{\vee}, \mathrm{p}_{1}\right) \vee \ldots \vee \operatorname{Card}\left(\mathrm{A} \boldsymbol{\bullet}, \mathrm{p}_{52}\right) \\
& \operatorname{Card}\left(K \vee, p_{1}\right) \vee \ldots \vee \operatorname{Card}\left(K \vee, p_{52}\right) \\
& \neg \operatorname{Card}\left(A \vee, p_{1}\right) \vee \neg \operatorname{Card}\left(A \vee, p_{2}\right) \\
& \neg \operatorname{Card}\left(A \vee, p_{1}\right) \vee \neg \operatorname{Card}\left(A \vee, p_{3}\right) \\
& \text { What will } \\
& \text { happen? }
\end{aligned}
$$

## Deck of Cards Graphically


[Van den Broeck 2015]

## Deck of Cards Graphically


[Van den Broeck 2015]

## Deck of Cards Graphically



One model/perfect matching
[Van den Broeck 2015]

## Deck of Cards Graphically


[Van den Broeck 2015]

## Deck of Cards Graphically


[Van den Broeck 2015]

## Deck of Cards Graphically



Model counting: How many perfect matchings?

## Deck of Cards Graphically


[Van den Broeck 2015]

## Deck of Cards Graphically



What if I add the unit clause $\neg \operatorname{Card}\left(\mathrm{K} \bullet, \mathrm{p}_{52}\right)$ to my CNF?

## Deck of Cards Graphically



What if I add the unit clause
$\neg \operatorname{Card}\left(\mathrm{K} \bullet, \mathrm{p}_{52}\right)$ to my CNF?

## Deck of Cards Graphically



What if I add unit clauses to my CNF?

## Observations

- Deck of cards = complete bigraph
- Unit clause removes edge

Encode any bigraph

- Counting models = perfect matchings
- Problem is \#P-complete! :
- All solvers efficiently handle unit clauses
- No solver can do cards problem efficiently!


## What's Going On Here?



Probability that Card52 is Spades given that Card1 is QH?

## What's Going On Here?



Probability that Card52 is Spades given that Card1 is QH?

13/51

## What's Going On Here?



Probability that Card52 is Spades given that Card2 is QH?

## What's Going On Here?



## Probability that Card52 is Spades given that Card2 is QH? <br> 13/51

## What's Going On Here?



Probability that Card52 is Spades given that Card3 is QH?

## What's Going On Here?



Probability that Card52 is Spades given that Card3 is QH?

13/51

## Tractable Reasoning



## What's going on here? <br> Which property makes reasoning tractable?

## Tractable Reasoning



## What's going on here?

Which property makes reasoning tractable?

- High-level (first-order) reasoning
- Symmetry
- Exchangeability

$\Rightarrow$ Lifted Inference

## What are first-order circuit languages?

## Negation Normal Form


[Darwiche 2002]

## Decomposable NNF



## Deterministic Decomposable NNF


[Darwiche 2002]

## Deterministic Decomposable NNF

Weighted Model Counting

[Darwiche 2002]

## Deterministic Decomposable NNF

Weighted Model Counting and much more!

[Darwiche 2002]

## First-Order NNF

## $\forall X, X \in \operatorname{People}: \operatorname{belgian}(X) \Rightarrow$ likes $(X$, chocolate $)$


[Van den Broeck 2013]

## First-Order Decomposability

$\forall X, X \in$ People : belgian $(X) \Rightarrow$ likes $(X$, chocolate $)$

## Decomposable


[Van den Broeck 2013]

## First-Order Decomposability

$$
\forall X, X \in \text { People }: \text { belgian }(X) \Rightarrow \text { likes }(X, \text { chocolate })
$$


[Van den Broeck 2013]

## First-Order Determinism

$\forall X, X \in \operatorname{People}: \operatorname{belgian}(X) \Rightarrow$ likes $(X$, chocolate $)$

[Van den Broeck 2013]

## Deterministic Decomposable FO NNF

$$
\forall X, X \in \text { People : belgian }(X) \Rightarrow \text { likes }(X, \text { chocolate })
$$

Weighted Model Counting

[Van den Broeck 2013]

## Deterministic Decomposable FO NNF

$$
\forall X, X \in \text { People : belgian }(X) \Rightarrow \text { likes }(X, \text { chocolate })
$$

Weighted Model Counting

[Van den Broeck 2013]

## Deterministic Decomposable FO NNF

$$
\forall X, X \in \text { People : belgian }(X) \Rightarrow \text { likes }(X, \text { chocolate })
$$

Weighted Model Counting

[Van den Broeck 2013]

## How to do first-order knowledge compilation?

## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(\mathrm{y}))
$$



## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$



## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(\mathrm{y}))
$$



## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$



## Deterministic

[Van den Broeck 2013]

## Deterministic Decomposable FO NNF

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(\mathrm{y}))
$$


[Van den Broeck 2013]

## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\operatorname{Smokes}(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

- If we know $\mathbf{D}$ precisely: who smokes, and there are $k$ smokers?


## Database:

$$
\begin{aligned}
& \text { Smokes(Alice) = } 1 \\
& \text { Smokes(Bob) = } 0 \\
& \text { Smokes(Charlie) = } 0 \\
& \text { Smokes(Dave) = } 1 \\
& \text { Smokes(Eve) }=0
\end{aligned}
$$

Smokes


Smokes


## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
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& \text { Smokes(Dave) = } 1 \\
& \text { Smokes(Eve) }=0
\end{aligned}
$$



## First-Order Model Counting: Example

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\end{aligned}
$$



## First-Order Model Counting: Example

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& \text { Smokes(Eve) }=0
\end{aligned}
$$



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& \text { Smokes(Dave) = } 1 \\
& \text { Smokes(Eve) }=0
\end{aligned}
$$



## First-Order Model Counting: Example

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\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
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& \text { Smokes(Dave) = } 1 \\
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\end{aligned}
$$



## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

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$$
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& \text { Smokes(Charlie) = } 0 \\
& \text { Smokes(Dave) = } \\
& \text { Smokes(Eve) = } 0
\end{aligned}
$$



## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

- If we know D precisely: who smokes, and there are $k$ smokers?


## Database:

$$
\begin{aligned}
& \text { Smokes(Alice) = } 1 \\
& \text { Smokes(Bob) = } 0 \\
& \text { Smokes(Charlie) = } 0 \\
& \text { Smokes(Dave) = } 1 \\
& \text { Smokes(Eve) }=0
\end{aligned}
$$



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& \text { Smokes(Charlie) }=0 \\
& \text { Smokes(Dave) }=1 \\
& \text { Smokes(Eve) }=0 \\
& \cdots \\
& \cdots \\
& 2^{n^{2}-k(n-k)} \text { models }
\end{aligned}
$$

Smokes


## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
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- If we know D precisely: who smokes, and there are $k$ smokers?

$$
\begin{aligned}
& \text { Database: } \\
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& \text { Smokes(Charlie) }=0 \\
& \text { Smokes(Dave) }=1 \\
& \text { Smokes(Eve) }=0 \\
& \cdots \\
& \cdots 2^{n^{2}-k(n-k) \quad \text { models }}
\end{aligned}
$$



- If we know that there are $k$ smokers?


## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

- If we know $\mathbf{D}$ precisely: who smokes, and there are $k$ smokers?


## Database:

Smokes(Alice) = 1
Smokes(Bob) $=0$
Smokes(Charlie) $=0$
Smokes(Dave) = 1
Smokes(Eve) $=0$

$$
\rightarrow 2^{n^{2}-k(n-k)} \text { models }
$$

- If we know that there are $k$ smokers?


$$
\rightarrow\binom{n}{k} 2^{n^{2}-k(n-k)} \quad \text { models }
$$

## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

- If we know $\mathbf{D}$ precisely: who smokes, and there are $k$ smokers?


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Smokes(Eve) $=0$

$$
\rightarrow 2^{n^{2}-k(n-k)} \text { models }
$$



- If we know that there are $k$ smokers?
- In total...


## First-Order Model Counting: Example

$$
\Delta=\forall x, y \in \text { People, }(\text { Smokes }(x) \wedge \text { Friends }(x, y) \Rightarrow \text { Smokes }(y))
$$

- If we know $\mathbf{D}$ precisely: who smokes, and there are $k$ smokers?


## Database:

$$
\begin{aligned}
& \text { Smokes(Alice) }=1 \\
& \text { Smokes(Bob) }=0 \\
& \text { Smokes(Charlie) }=0 \\
& \text { Smokes(Dave) }=1 \\
& \text { Smokes(Eve) }=0 \\
& \ldots \\
& \rightarrow 2^{n^{2}-k(n-k)} \text { models }
\end{aligned}
$$



- If we know that there are $k$ smokers? $\quad \rightarrow\binom{n}{k} 2^{n^{2}-k(n-k)}$ models
- In total...

$$
\rightarrow \quad \sum_{k=0}^{n}\binom{n}{k} 2^{n^{2}-k(n-k)} \text { models }
$$

## Compilation Rules

- Standard rules
- Shannon decomposition (DPLL)
- Detect decomposability
- Etc.
- FO Shannon decomposition:



## Playing Cards Revisited

## Let us automate this:



$$
\begin{array}{r}
\forall p, \exists c, \operatorname{Card}(p, c) \\
\forall c, \exists p, \operatorname{Card}(\mathrm{p}, \mathrm{c}) \\
\forall \mathrm{p}, \forall \mathrm{c}, \forall \mathrm{c}^{\prime}, \operatorname{Card}(\mathrm{p}, \mathrm{c}) \wedge \operatorname{Card}\left(\mathrm{p}, \mathrm{c}^{\prime}\right) \Rightarrow \mathrm{c}=\mathrm{c}^{\prime}
\end{array}
$$

## Playing Cards Revisited

## Let us automate this:



$$
\text { \#SAT }=\sum_{k=0}^{n}\binom{n}{k} \sum_{l=0}^{n}\binom{n}{l}(l+1)^{k}(-1)^{2 n-k-l}=\mathrm{n}!
$$

## Playing Cards Revisited

## Let us automate this:



$$
\text { \#SAT }=\sum_{k=0}^{n}\binom{n}{k} \sum_{l=0}^{n}\binom{n}{l}(l+1)^{k}(-1)^{2 n-k-l}=\mathrm{n}!
$$

## Computed in time polynomial in $n$

## Perspectives...

## What I did not talk about... in KC

- Other queries and transformations
(see Dan Olteanu poster)
- Other KC languages
(FO-AODD)
- KC for logic programs
(see Vlasselaer poster)


## What I did not talk about...in FOMC

- WFOMC for probabilistic databases (see Gribkoff poster)
- WFOMC for probabilistic programs (see Vlasselaer poster)
- Complexity theory (data or domain)
- PTime domain complexity for 2-var fragment
- \# $\mathrm{P}_{1}$ domain complexity for some 3-var CNFs


## What I did not talk about...in FO

- Very related problems
- Lifted inference in SRL
- Very related applications
- Approximate lifted inference in Markov Logic
- Learn Markov logic networks
- Classical first-order reasoning
- Answer set programming,
- SMT,
- Theorem proving


## Format for First-Order BeyondNP

- DIMACS contributed to SAT success
- Goals
- Trivial to parse
- Captures MLNs, Prob. Programs, Prob. DBs
- Not a powerful representation language
- FO-CNF format under construction
- Vibhav?

```
p fo-cnf 2 1
d people 1000
r Friends(people,people)
r Smokes(people)
-Smokes(x) -Friends(x,y) Smokes(y)
w Friends 3.5 1.2
w Smokes -0.5 4
```


## Calendar

## At IJCAI in New York on July 9-11



- StarAI 2016 (http://www.starai. org/2016) Sixth International Workshop on Statistical Relational AI
- IJCAI Tutorial
"Lifted Probabilistic Inference in Relational Models" with Dan Suciu


## Conclusions

- FOMC is BeyondNP reduction target
- Existing solvers inadequate

Exponential speedups from FO solvers

- FOKC is Elegant, more than FOMC
- Intersection of communities
- Statistical relational learning (lifted inference)
- Probabilistic databases
- Automated reasoning (you!)


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