iscc Tutorial

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Outline

1. Introduction
2. Basic Concepts and Operations
    - Sets and Iteration Domains
    - Maps and Code Generation
    - Access Relations and Polyhedral Model
    - Dependence Analysis
    - Transitive Closures
    - Basic Counting
    - Computing Bounds
    - Weighted Counting
3. Simple Applications
    - Pointer Conversion
    - Dynamic Memory Requirement Estimation
    - Reuse Distance Computation

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  - ⇒ interactive interface to the `barvinok` counting library
  - ⇒ also provides interface to the `CLooG` code generation library, to the `pet` polyhedral model extractor and to some operations of the `isl` integer set library
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  - ⇒ available from `http://freecode.com/projects/barvinok/`

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- How to run `iscc`?
  - ⇒ compile and install `barvinok` following the instructions in `README`
  - ⇒ run `iscc`
    Note: `iscc` currently does not use readline, so you may want to use a readline front-end: `rlwrap iscc`

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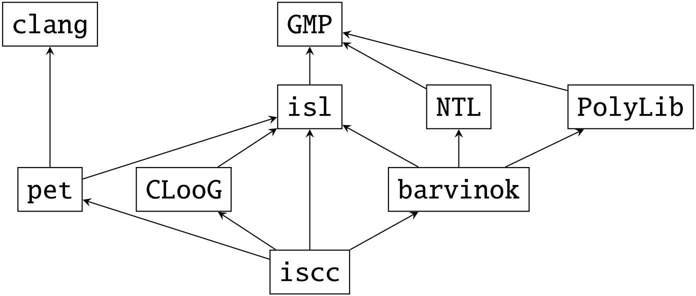
Interaction with Libraries

`isl`: manipulates parametric affine sets and relations

`barvinok`: counts elements in parametric affine sets and relations

`CLooG`: generates code to scan elements in parametric affine sets

`pet`: extracts polyhedral model



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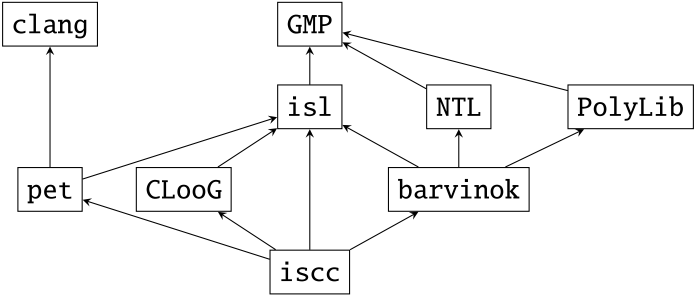

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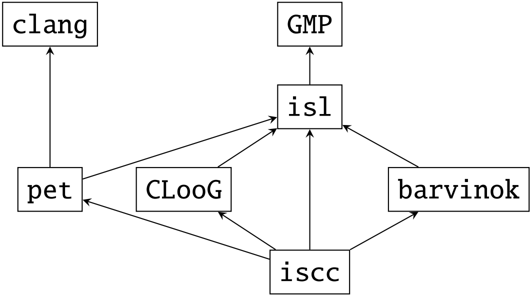

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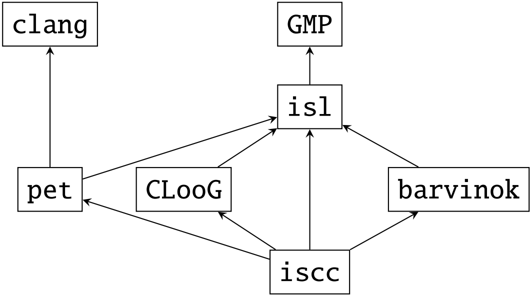

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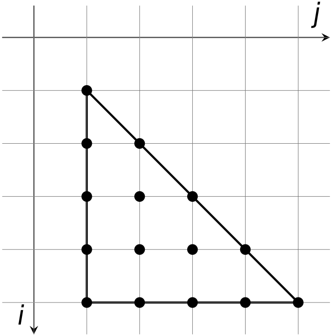

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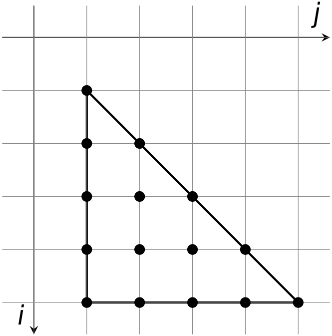

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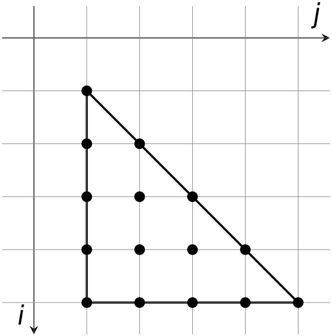

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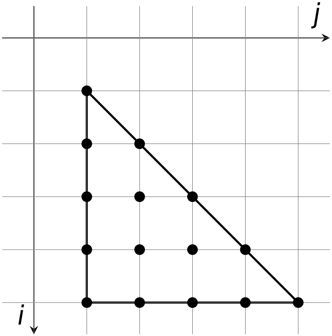

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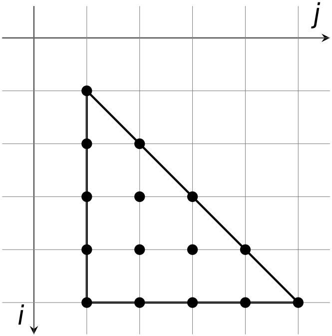

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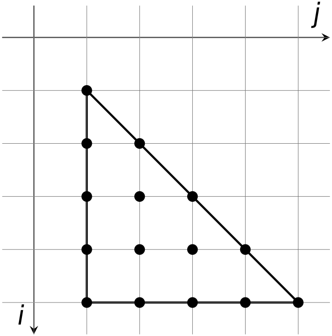

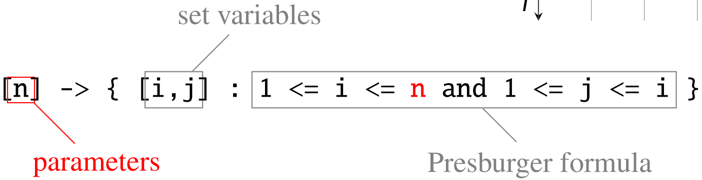

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- set variables
    - local to set
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```

is equal to

```
[n] -> { [a,b] : 1 <= a <= n and 1 <= b <= a }
```

but not equal to

```
[n] -> { [j,i] : 1 <= i <= n and 1 <= j <= i }
```

or

```
[m] -> { [i,j] : 1 <= i <= m and 1 <= j <= i }
```

Code Generation, Schedules and Maps

```
for (i = 1; i <= n; ++i)
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⇒ generate code that visits elements in lexicographic order

`codegen1`, `codegen2`

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What if a different order is needed?
⇒ apply a schedule: maps iterations domain to multi-dimensional time
⇒ multi-dimensional time is ordered lexicographically

Example: interchange `i` and `j`

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intersect domain of map on the left with set on the right

codegen1, codegen2

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Access Relations and Polyhedral Model

Simple program with temporary array `t`:

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for (i = 0; i < N; ++i)
S1:     t[i] = f(a[i]);
for (i = 0; i < N; ++i)
S2:     b[i] = g(t[N-i-1]);
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An access relation maps an iteration to an array index

For example, the access relation for the read in `S2`:

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Polyhedral model of a program consists of

- iteration domains
- access relations (reads and writes)
- schedule

```
M := parse_file("simple.c");
D := M[0]; W := M[1]; R := M[2]; S := M[3];
```

Lexicographic Optimization

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
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- What is the last iteration of the loop?

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Dependence Analysis

> *Given a read from an array element, what was the last write to the same array element before the read?*

Simple case: array written through a single access

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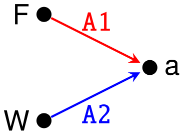

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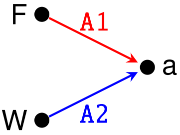

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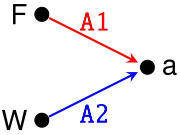

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In general: impose lexicographical order on shared iterators

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`last Write before Read under Schedule`

Result: last write + set of reads without corresponding write

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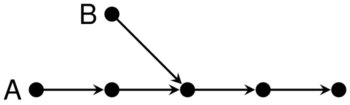
F:        t[j] = f(t[j], t[j+1]);
for (i = 0; i < n; ++i)
B:   b[i] = t[i];
```

```
M := parse_file("dep.c");
Write := M[1]; Read := M[2]; Sched := M[3];
last Write before Read under Sched;
```

Transitive Closures

Given a graph (represented as an affine map)

```
M := { A[i] -> A[i+1] : 0 <= i <= 3; B[] -> A[2] };
```



What is the transitive closure?

Transitive Closures

Given a graph (represented as an affine map)

```
M := { A[i] -> A[i+1] : 0 <= i <= 3; B[] -> A[2] };
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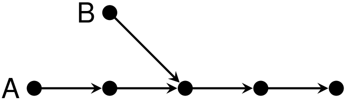
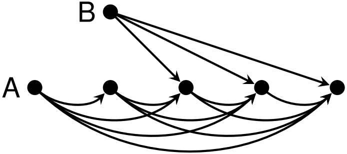


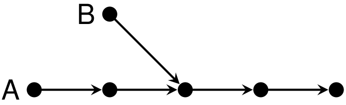
What is the transitive closure? ⇒ `M^+;`



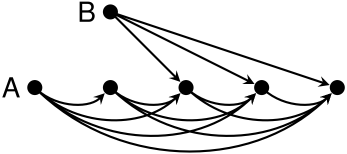
Transitive Closures

Given a graph (represented as an affine map)

```
M := { A[i] -> A[i+1] : 0 <= i <= 3; B[] -> A[2] };
```



What is the transitive closure? ⇒ `M^+;`



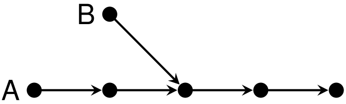
Result:

```
({ B[] -> A[o0] : o0 <= 4 and o0 >= 3; B[] -> A[2];
   A[i] -> A[o0] : i >= 0 and i <= 3 and o0 >= 1 and
   o0 <= 4 and o0 >= 1 + i }, True)
```

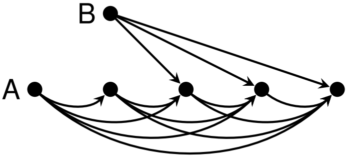
Transitive Closures

Given a graph (represented as an affine map)

```
M := { A[i] -> A[i+1] : 0 <= i <= 3; B[] -> A[2] };
```



What is the transitive closure? ⇒ `M^+;`



Result:

exact transitive closure

```
({ B[] -> A[o0] : o0 <= 4 and o0 >= 3; B[] -> A[2];
   A[i] -> A[o0] : i >= 0 and i <= 3 and o0 >= 1 and
   o0 <= 4 and o0 >= 1 + i }, True)
```

Reachability Analysis

```
double x[2][10];
int old = 0, new = 1, i, t;
for (t = 0; t<1000; t++) {
  for (i = 0; i<10;i++)
    x[new][i] = g(x[old][i]);
  new = (new+1) %2; old = (old+1) %2;
}
```

Invariant between `new` and `old`?

Reachability Analysis

```
double x[2][10];
int old = 0, new = 1, i, t;
for (t = 0; t<1000; t++) {
  for (i = 0; i<10;i++)
    x[new][i] = g(x[old][i]);
  new = (new+1) %2; old = (old+1) %2;
}
```

Invariant between `new` and `old`?

```
T := {[new,old] -> [(new+1)%2,(old+1)%2]};
S0 := {[0,1]};
(T^+)(S0);
```

Cardinality

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

- How many times is the statement executed?

```
S := [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
card S;
```

Cardinality

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

- How many times is the statement executed?

```
S := [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
card S;
```

number of elements in the set

Cardinality

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

- How many times is the statement executed?

```
S := [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
card S;
```

number of elements in the set

- How many times is a given array element written?

```
A:=[N]->{[i,j]->a[i+j]:0<=i<N and 0<=j<N-i};
card (A^-1);
```

Cardinality

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

- How many times is the statement executed?

```
S := [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
card S;
```

number of elements in the set

- How many times is a given array element written?

```
A:=[N]->{[i,j]->a[i+j]:0<=i<N and 0<=j<N-i};
card (A^-1);
```

number of image elements

Cardinality

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

- How many times is the statement executed?

```
S := [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
card S;
```

number of elements in the set

- How many times is a given array element written?

```
A:=[N]->{[i,j]->a[i+j]:0<=i<N and 0<=j<N-i};
card (A^-1);
```

number of image elements

- How many array elements are written?

```
A:=[N]->{[i,j]->a[i+j]:0<=i<N and 0<=j<N-i};
card (ran A);
```

`card1,card2,card3`

Quasipolynomials

```
for (i = 1; i <= n; ++i)
    for (j = 1; j <= n - 2 * i; ++j)
        /* S */
```

How many times is S executed?

```
card [n] -> { [i,j] : 1 <= i <= n and 1 <= j <= n - 2i };
```

Quasipolynomials

```
for (i = 1; i <= n; ++i)
    for (j = 1; j <= n - 2 * i; ++j)
        /* S */
```

How many times is S executed?

```
card [n] -> { [i,j] : 1 <= i <= n and 1 <= j <= n - 2i };
```

Result:

```
[n] -> { ((-1/4 * n + 1/4 * n^2) - 1/2 * [(n)/2]) : n >= 3 }
```

That is,

$$-\frac{n}{4} + \frac{n^2}{4} - \frac{1}{2}\left\lfloor \frac{n}{2} \right\rfloor \quad \text{if } n \geq 3.$$

Quasipolynomials

```
for (i = 1; i <= n; ++i)
    for (j = 1; j <= n - 2 * i; ++j)
        /* S */
```

How many times is S executed?

```
card [n] -> { [i,j] : 1 <= i <= n and 1 <= j <= n - 2i };
```

Result:

greatest integer part

```
[n] -> { ((-1/4 * n + 1/4 * n^2) - 1/2 * [(n)/2]) : n >= 3 }
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Quasipolynomials

```
for (i = 1; i <= n; ++i)
    for (j = 1; j <= n - 2 * i; ++j)
        /* S */
```

How many times is S executed?

```
card [n] -> { [i,j] : 1 <= i <= n and 1 <= j <= n - 2i };
```

greatest integer part

Result:

```
[n] -> { ((-1/4 * n + 1/4 * n^2) - 1/2 * [(n)/2]) : n >= 3 }
```

That is,

$$-\frac{n}{4}+\frac{n^2}{4}-\frac{1}{2}\left\lfloor\frac{n}{2}\right\rfloor \quad \text{if } n \geq 3.$$

Polynomial approximations

⇒ run `iscc --polynomial-approximation`

Memory Requirements

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p = malloc(i * j + i - N + 1);
        /* ... */
        free(p);
    }
```

How much memory is needed?

Memory Requirements

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p = malloc(i * j + i - N + 1);
        /* ... */
        free(p);
    }
```

How much memory is needed?

```
ub [N] -> {[i,j] -> i*j+i-N+1: 0 <= i < N and i <= j < N};
```

Memory Requirements

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p = malloc(i * j + i - N + 1);
        /* ... */
        free(p);
    }
```

How much memory is needed?

```
ub [N] -> {[i,j] -> i*j+i-N+1: 0 <= i < N and i <= j < N};
```

Result:

```
([N] -> { max((1 - 2 * N + N^2)) : N >= 1 }, True)
```

Memory Requirements

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p = malloc(i * j + i - N + 1);
        /* ... */
        free(p);
    }
```

How much memory is needed?

```
ub [N] -> {[i,j] -> i*j+i-N+1: 0 <= i < N and i <= j < N};
```

Result:

```
([N] -> { max((1 - 2 * N + N^2)) : N >= 1 }, True)
```

bound is tight

Incremental Counting

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

How many times is the statement executed?

- direct computation

```
card [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };
```

Incremental Counting

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

How many times is the statement executed?

- direct computation

  `card [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };`

- incremental computation

  `card [N] -> { [i] -> [j] : 0<=i<N and 0<=j<N-i };`

Incremental Counting

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

How many times is the statement executed?

- direct computation

  `card [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };`

- incremental computation

  `card [N] -> { [i] -> [j] : 0<=i<N and 0<=j<N-i };`

  Result:

  ```
  [N] -> { [i] -> (N - i) : i <= -1 + N and i >= 0 }
  ```

  ```
  sum [N] -> { [i] -> (N - i) : i <= -1 + N and i >= 0 };
  ```

Incremental Counting

```
for (i = 0; i < N; ++i)
    for (j = 0; j < N - i; ++j)
        a[i+j] = f(a[i+j]);
```

How many times is the statement executed?

- direct computation

  `card [N] -> { [i,j] : 0<=i<N and 0<=j<N-i };`

- incremental computation

  `card [N] -> { [i] -> [j] : 0<=i<N and 0<=j<N-i };`

  Result:
  `[N] -> { [i] -> (N - i) : i <= -1 + N and i >= 0 }`

  `sum [N] -> { [i] -> (N - i) : i <= -1 + N and i >= 0 };`

  sum over all elements in domain

Total Memory Allocation

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j)
        p[i][j] = malloc(i * j + i - N + 1);
/* ... */
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j)
        free(p[i][j]);
```

How much memory allocated in total?

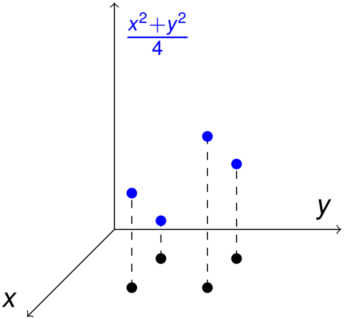
Total Memory Allocation

```
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j)
        p[i][j] = malloc(i * j + i - N + 1);
/* ... */
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j)
        free(p[i][j]);
```

How much memory allocated in total?

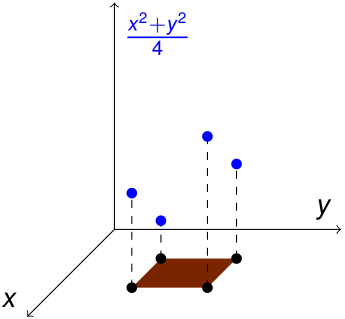
```
sum [N] -> {[i,j] -> i*j+i-N+1: 0 <= i < N and i <= j < N};
```

Weighted Counting



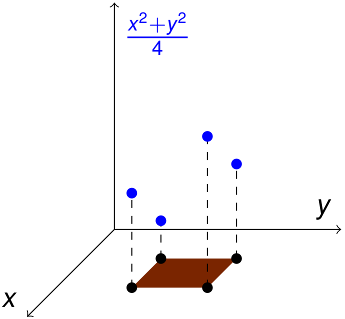
```
F := { [x,y] -> 1/4*x^2+1/4*y^2 : 1 <= x,y <= 2 };
```

Weighted Counting



```
F := { [x,y] -> 1/4*x^2+1/4*y^2 : 1 <= x,y <= 2 };
D := dom F;
```

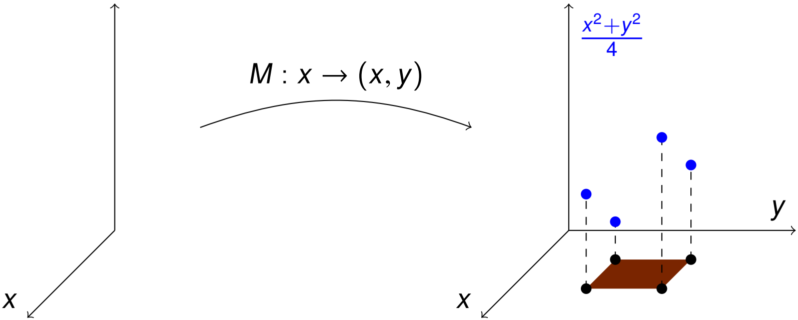
Weighted Counting



```
F := { [x,y] -> 1/4*x^2+1/4*y^2 : 1 <= x,y <= 2 };
D := dom F;
F(D);
```

⇒ sum of `F` over points in `D`

Weighted Counting



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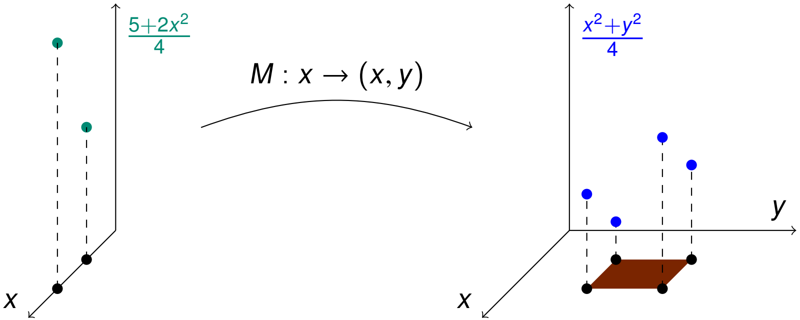
F(D);
```

⇒ sum of `F` over points in `D`

```
M := { [x] -> [x,y] };
```

Weighted Counting



```
F := { [x,y] -> 1/4*x^2+1/4*y^2 : 1 <= x,y <= 2 };
D := dom F;
F(D);
```

⇒ sum of `F` over points in `D`

```
M := { [x] -> [x,y] };
F(M);
```

⇒ sum of `F` over image of `M` (alternative notation: `M . F`)

Compositions with Piecewise (Folds of) Quasipolynomials

`f . g;`

- `f`: $D_1 \to D_2$ is a map
- `g`: $D_2 \to \mathbb{Q}$ may be
  - piecewise quasipolynomial
    (result of counting problems)

    $\Rightarrow$ take sum over intersection of `ran f` and `dom g`
  - piecewise fold of quasipolynomials
    (result of upper bound computation)

    $\Rightarrow$ compute bound over intersection of `ran f` and `dom g`
- (`f . g`): $D_1 \to \mathbb{Q}$ of same type as `g`

Note: if f is single-valued, then sum/bound is computed over a single point

Outline

1. Introduction
2. Basic Concepts and Operations
    - Sets and Iteration Domains
    - Maps and Code Generation
    - Access Relations and Polyhedral Model
    - Dependence Analysis
    - Transitive Closures
    - Basic Counting
    - Computing Bounds
    - Weighted Counting
3. Simple Applications
    - Pointer Conversion
    - Dynamic Memory Requirement Estimation
    - Reuse Distance Computation

Pointer Conversion

```
p = a;
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p += j * ((j-i)/4);
        *p = hard_work(i,j);
    }
```

Can we parallelize this code?

Pointer Conversion

```
p = a;
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
        p += j * ((j-i)/4);
        *p = hard_work(i,j);
    }
```

Can we parallelize this code?

⇒ No, (false) dependency through `p`

⇒ Compute closed formula for `p`

$$p = a + \sum_{\substack{(i',j') \in S \\ (i',j') \preccurlyeq (i,j)}} j' \left\lfloor \frac{j' - i'}{4} \right\rfloor$$

with $S = \{ (i',j') \in \mathbb{Z}^2 \mid 0 \leq i' < N \wedge i' \leq j' < N \}$

Pointer Conversion

```
p = a;
for (i = 0; i < N; ++i)
    for (j = i; j < N; ++j) {
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⇒ No, (false) dependency through `p`

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$$p = a + \sum_{\substack{(i',j') \in S \\ (i',j') \boxed{\preccurlyeq} (i,j)}} j' \left\lfloor \frac{j'-i'}{4} \right\rfloor$$

lexicographically less than

with $S = \{\, (i',j') \in \mathbb{Z}^2 \mid 0 \le i' < N \wedge i' \le j' < N \,\}$

Pointer Conversion

$$p = a + \sum_{\substack{(i',j')\in S \\ (i',j')\preccurlyeq(i,j)}} j' \left\lfloor \frac{j'-i'}{4} \right\rfloor$$

with $S = \{\, (i', j') \in \mathbb{Z}^2 \mid 0 \leq i' < N \wedge i' \leq j' < N \,\}$

Pointer Conversion

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```
S := [N] -> { [i,j] : 0 <= i < N and i <= j < N };
L := S <<= S;
INC := { [[i,j] -> [i',j']] -> j' * [(j'-i')/4] };
INC := INC * (wrap (L^-1));
sum INC;
```

Pointer Conversion

$$p = a + \sum_{\substack{(i',j')\in S \\ (i',j')\preccurlyeq(i,j)}} j' \left\lfloor \frac{j'-i'}{4} \right\rfloor$$

with $S = \{ (i',j') \in \mathbb{Z}^2 \mid 0 \le i' < N \wedge i' \le j' < N \}$

map: (elements of) left set lexicographically smaller than right set

```
S := [N] -> { [i,j] : 0 <= i < N and i <= j < N };
L := S <<= S;
INC := { [[i,j] -> [i',j']] -> j' * [(j'-i')/4] };
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```

embed map in a set

Pointer Conversion

$$p = a + \sum_{\substack{(i',j')\in S\\ (i',j')\preccurlyeq(i,j)}} j' \left\lfloor \frac{j'-i'}{4} \right\rfloor$$

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INC := { [[i,j] -> [i',j']] -> j' * [(j'-i')/4] };
INC := INC * (wrap (L^-1));
sum INC;
```

embed map in a set

Note: if domain of argument to `sum` [ub] is an embedded map, then sum [bound] is computed over range of embedded map

Dynamic Memory Requirement Estimation [CFGV2006]

How much memory is needed to execute the following program?

```
void m0(int m) {
  for (c = 0; c < m; c++) {
    m1(c);                      /*S1*/
    B[] m2Arr = m2(2*m-c);      /*S2*/
  }
}
void m1(int k) {
  for (i = 1; i <= k; i++) {
    A a = new A();              /*S3*/
    B[] dummyArr = m2(i);       /*S4*/
  }
}
B[] m2(int n) {
  B[] arrB = new B[n];          /*S5*/
  for (j = 1; j <= n; j++)
    B b = new B();              /*S6*/
  return arrB;
}
```

Dynamic Memory Requirement Estimation [CFGV2006]

How much memory is needed to execute the following program?

```
void m0(int m) {
  for (c = 0; c < m; c++) {
    m1(c);                    /*S1*/
    B[] m2Arr = m2(2*m-c);    /*S2*/
  }
}
void m1(int k) {
  for (i = 1; i <= k; i++) {
    A a = new A();            /*S3*/
    B[] dummyArr = m2(i);     /*S4*/
  }
}
B[] m2(int n) {
  B[] arrB = new B[n];        /*S5*/
  for (j = 1; j <= n; j++)
    B b = new B();            /*S6*/
  return arrB;
}
```

```
D := {
m0[m]->S1[c] : 0<=c<m;
m0[m]->S2[c] : 0<=c<m;
m1[k]->S3[i] : 1<=i<=k;
m1[k]->S4[i] : 1<=i<=k;
m2[n]->S5[];
m2[n]->S6[j] : 1<=j<=n
};
DM := (domain_map D)^-1;
```

Dynamic Memory Requirement Estimation [CFGV2006]

How much (scoped) memory is needed?
⇒ compute for each method

$\mathrm{ret}_{\mathtt{m}}$	size of memory returned by `m`
$\mathrm{cap}_{\mathtt{m}}$	size of memory “captured” (not returned) by `m`
$\mathrm{memRq}_{\mathtt{m}}$	total memory requirements of `m`

$$\mathrm{memRq}_{\mathtt{m}} = \mathrm{cap}_{\mathtt{m}} + \max_{\text{p called by } \mathtt{m}} \mathrm{memRq}_{\mathtt{p}}$$

Dynamic Memory Requirement Estimation [CFGV2006]

How much (scoped) memory is needed?
⇒ compute for each method

$\mathrm{ret}_{\mathtt{m}}$	size of memory returned by `m`
$\mathrm{cap}_{\mathtt{m}}$	size of memory “captured” (not returned) by `m`
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B[] m2(int n) {
 B[] arrB = new B[n];
 for (j=1; j<=n; j++)
  B b = new B();
 return arrB;
}
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}
```

```
ret_m2 := DM .
   { [m2[n] -> S5[]] -> n : n >= 0 };
cap_m2 := DM .
   { [m2[n] -> S6[j]] -> 1 };
req_m2 := cap_m2 +
   { m2[n] -> max(0) };
```

Dynamic Memory Requirement Estimation [CFGV2006]

```
void m1(int k) {
    for (i = 1; i <= k; i++) {
        A a = new A();              /* S3 */
        B[] dummyArr = m2(i);       /* S4 */
   }
}
```

$$\mathrm{cap}_{\mathtt{m1}}(k) = \sum_{1 \le i \le k} (1 + \mathrm{ret}_{\mathtt{m2}}(i))$$

`ret_m2` is a function of the arguments of `m2`

We want to use it as a function of the arguments and local variables of `m1`

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⇒ define parameter binding

```
CB_m1 := { [m1[k] -> S4[i]] -> m2[i] };
cap_m1 := DM . ({ [m1[k]->S3[i]] -> 1 } + (CB_m1 . ret_m2));
```

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void m1(int k) {
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    }
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```

$$\mathrm{memRq_m} = \mathrm{cap_m} + \max_{\text{p called by m}} \mathrm{memRq_p}$$

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CB_m1 := { [m1[k] -> S4[i]] -> m2[i] };
ret_m1 := { m1[k] -> 0 };
cap_m1 := DM . ({ [m1[k]->S3[i]] -> 1 } + (CB_m1 . ret_m2));
req_m1 := cap_m1 + (DM . CB_m1 . req_m2);
```

Dynamic Memory Requirement Estimation [CFGV2006]

```
void m0(int m) {
    for (c = 0; c < m; c++) {
        m1(c);                          /* S1 */
        B[] m2Arr = m2(2 * m - c); /* S2 */
    }
}
```

```
CB_m0  := { [m0[m] -> S1[c]] -> m1[c];
            [m0[m] -> S2[c]] -> m2[2 * m - c] };
ret_m0  := { m0[m] -> 0 };
cap_m0  := DM . CB_m0 . (ret_m1 + ret_m2);
req_m0  := cap_m0 + (DM . CB_m0 . (req_m1 . req_m2));
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```

combine reductions

Reuse Distance Computation

Given an access to a cache line ℓ, how many distinct cache lines have been accessed since the previous access to ℓ?
⇒ Is the cache line still in the cache?

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⇒ Is the cache line still in the cache?

```
for (i = 0; i <= 7; ++i) {
    A[i];               //reference a
    A[7-i];             //reference b
    if (i <= 3)
        A[2*i];         //reference c
}
```

Assume `A[i]` in cache line $\lfloor i/3 \rfloor$

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Assume `A[i]` in cache line $\lfloor i/3 \rfloor$

i	0			1			2			3			4		5		6		7	
r	a	b	c	a	b	c	a	b	c	a	b	c	a	b	a	b	a	b	a	b
$r@i$	0	7	0	1	6	2	2	5	4	3	4	6	4	3	5	2	6	1	7	0
$\lfloor (r@i)/3 \rfloor$	0	2	0	0	2	0	0	1	1	1	1	2	1	1	1	0	2	0	2	0
distance	0	0	2	1	2	2	1	0	1	1	1	3	2	1	1	3	3	2	2	2

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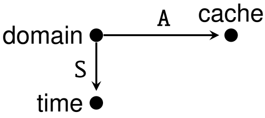

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```
D := { a[i] : 0 <= i <= 7; b[i] : 0 <= i <= 7; c[i] : 0 <= i <= 3 };
C := { A[i] -> L[j] : exists a = [i/3] : j = a };
A := ({ a[i] -> A[i]; b[i] -> A[7-i]; c[i] -> A[2i] } . C) * D;
S := { a[i] -> [i,0]; b[i] -> [i,1]; c[i] -> [i,2] } * D;
```

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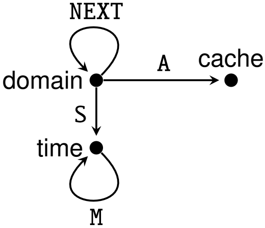

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A := ({ a[i] -> A[i]; b[i] -> A[7-i]; c[i] -> A[2i] } . C) * D;
S := { a[i] -> [i,0]; b[i] -> [i,1]; c[i] -> [i,2] } * D;
TIME := ran S; LT := TIME << TIME; LE := TIME <<= TIME;
T := ((S^-1) . A . (A^-1) . S) * LT;
M := lexmin T;
NEXT := S . M . (S^-1); # map to next access to same cache line
AFTER_PREV := (NEXT^-1) . (S . LE . (S^-1));
BEFORE := S . (LE^-1) . (S^-1);
card ((AFTER_PREV * BEFORE) . A);
```