

## From File Systems to Databases

Many Slides from Chapter 1 of  
**Database Systems: Design, Implementation, and  
Management, Fifth Edition, Rob and Coronel**

## Importance of DBMS

- **Makes data management more efficient and effective**
- **Query language allows quick answers to *ad hoc* queries**
- **Provides better access to more and better-managed data**
- **Promotes integrated view of organization's operations**
- **Reduces the probability of inconsistent data**

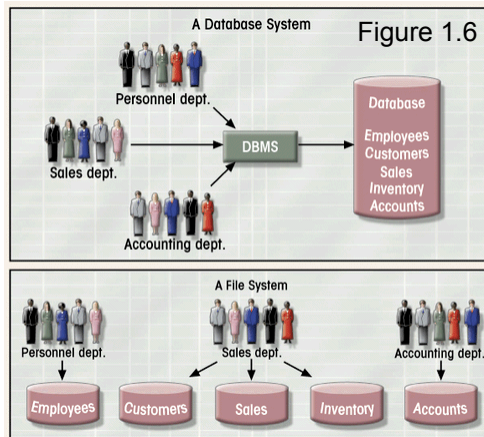
## Historical Roots of Database

- **First applications focused on clerical tasks**
- **Requests for information quickly followed**
- **File systems developed to address needs**
  - Data organized according to expected use
  - Data Processing (DP) specialists computerized manual file systems

## Database Systems

- **Database consists of logically related data stored in a single repository**
- **Provides advantages over file system management approach**
  - Eliminates inconsistency, data anomalies, data dependency, and structural dependency problems
  - Stores data structures, relationships, and access paths

## Database vs. File Systems



## Database System Types

- **Single-user vs. Multiuser Database**
  - Desktop
  - Workgroup
  - Enterprise
- **Centralized vs. Distributed**
- **Use**
  - Production or transactional
  - Decision support or data warehouse

## DBMS Functions

- Data dictionary management
- Data storage management
- Data transformation and presentation
- Security management
- Multiuser access control
- Backup and recovery management
- Data integrity management
- Database language and application programming interfaces
- Database communication interfaces

## Database Models

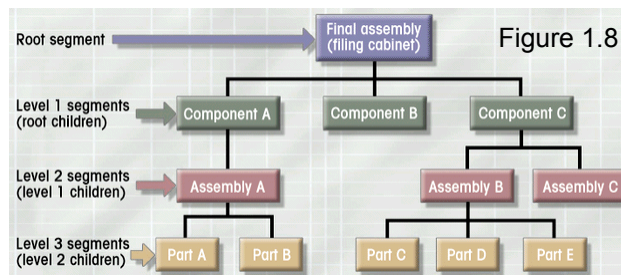
- Collection of logical constructs used to represent data structure and relationships within the database
  - Conceptual models: logical nature of data representation
  - Implementation models: emphasis on how the data are represented in the database

## Database Models (con't.)

- **Relationships in Conceptual Models**
  - One-to-one (1:1)
  - One-to-many (1:M)
  - Many-to-many (M:N)
- **Implementation Database Models**
  - Hierarchical
  - Network
  - Relational

## Hierarchical Database Model

- **Logically represented by an upside down tree**
  - Each parent can have many children
  - Each child has only one parent



## Network Database Model

- Each record can have multiple parents
  - Composed of sets
  - Each set has owner record and member record
  - Member may have several owners

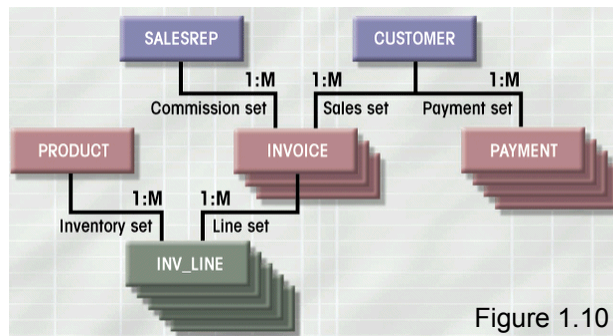


Figure 1.10

## Relational Database Model (70s)

- Tables are a series of row/column intersections
- Logical Model ... with great mathematical foundations—initially expected to be very inefficient
- Twenty years of R&D changed that dramatically: e.g. B+-trees & Optimizers have produced
  - Superior Performance with Data Independence
- Others Pluses: Standards & Strong vendors make RDBMS unbeatable
- Alternatives have resulted in extensions and improvements, rather than replacements.

## Relational DBMS Evolution

- **ER model**
- **Recursion**
- **O-O**
- **OLAP and decision support**
- **The web and XML, ...**

## Entity Relationship Database Model (Peter Chen '76)

- **Complements the relational data model concepts**
- **Represented in an entity relationship diagram (ERD)**
- **Based on entities, attributes, and relationships**

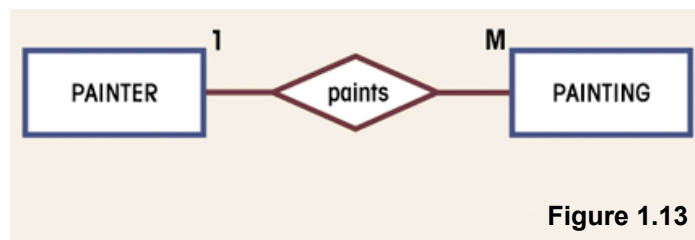


Figure 1.13

## Entity Relationship Database Model

- **Advantages**
  - Exceptional conceptual simplicity
  - Visual representation
  - Effective communication tool
  - Integrated with the relational database model
- **Disadvantages**
  - No data manipulation language
- **ER Model became a powerful design tool in DB design!**



## Recursion in SQL (80s)

- SQL:1999 permits recursive view definition
- Example: find all employee-manager pairs, where the employee reports to the manager directly or indirectly (that is manager's manager, manager's manager's manager, etc.)

```
with recursive empl (employee_name, manager_name) as (
    select employee_name, manager_name
    from   manager
union
    select manager.employee_name, empl.manager_name
    from   manager, empl
    where  manager.manager_name = empl.employee_name)
select *
from     empl
```

This example view, *empl*, is called the *transitive closure* of the *manager* relation







## The Power of Recursion

- Recursive views make it possible to write queries, such as transitive closure queries, that cannot be written without recursion or iteration.
  - Intuition: Without recursion, a non-recursive non-iterative program can perform only a fixed number of joins of *manager* with itself
    - ▶ This can give only a fixed number of levels of managers
    - ▶ Given a program we can construct a database with a greater number of levels of managers on which the program will not work
- Computing transitive closure
  - The next slide shows a *manager* relation
  - Each step of the iterative process constructs an extended version of *empl* from its recursive definition.
  - The final result is called the *fixed point* of the recursive view definition.
- Recursive views are required to be *monotonic*. That is, if we add tuples to *manager* the view contains all of the tuples it contained before, plus possibly more



## Example of Fixed-Point Computation

<i>employee_name</i>	<i>manager_name</i>
Alon	Barinsky
Barinsky	Estovar
Corbin	Duarte
Duarte	Jones
Estovar	Jones
Jones	Klinger
Rensal	Klinger

<i>Iteration number</i>	<i>Tuples in empl</i>
0	
1	(Duarte), (Estovar)
2	(Duarte), (Estovar), (Barinsky), (Corbin)
3	(Duarte), (Estovar), (Barinsky), (Corbin), (Alon)
4	(Duarte), (Estovar), (Barinsky), (Corbin), (Alon)





## Object-Relational Data Models

- Small Talk and C++ made OO very popular (90s)
- OODBMS featuring persistent types gained some popularity
- Vendors Extended the relational data model by including object orientation and constructs to deal with added data types (ORDBMS--SQL 2003).
- ORDBMS:
  - Allow attributes of tuples to have complex types, including non-atomic values such as nested relations.
  - Preserve relational foundations, in particular the declarative access to data, while extending modeling power.
  - Upward compatibility with existing relational languages.



## Complex Data Types

- Motivation:
  - Permit non-atomic domains (atomic = indivisible)
  - Example of non-atomic domain: set of integers, or set of tuples
  - Allows more intuitive modeling for applications with complex data
- Intuitive definition:
  - allow relations whenever we allow atomic (scalar) values
    - relations within relations
  - Retains mathematical foundation of relational model
  - Violates first normal form.





## Structured Types and Inheritance in SQL

- **Structured types** can be declared and used in SQL

```
create type Name as
  (firstname  varchar(20),
   lastname   varchar(20))
final
```

```
create type Address as
  (street     varchar(20),
   city       varchar(20),
   zipcode    varchar(20))
not final
```

- Note: **final** and **not final** indicate whether subtypes can be created
- Structured types can be used to create tables with composite attributes

```
create table customer (
  name      Name,
  address   Address,
  dateOfBirth date)
```
- Dot notation used to reference components: *name.firstname*



## Structured Types (cont.)

- User-defined row types

```
create type CustomerType as (
  name Name,
  address Address,
  dateOfBirth date)
```

- Can then create a table whose rows are a user-defined type

```
create table customer of CustomerType
```





## Methods

- Can add a method declaration with a structured type.

```
method ageOnDate (onDate date)
```

```
  returns interval year
```

- Method body is given separately.

```
create instance method ageOnDate (onDate date)
```

```
  returns interval year
```

```
  for CustomerType
```

```
  begin
```

```
    return onDate - self.dateOfBirth;
```

```
  end
```

- We can now find the age of each customer:

```
select name.lastname, ageOnDate (current_date)
```

```
from customer
```



## Inheritance

- Suppose that we have the following type definition for people:

```
create type Person
```

```
  (name varchar(20),
```

```
   address varchar(20))
```

- Using inheritance to define the student and teacher types

```
create type Student
```

```
  under Person
```

```
  (degree varchar(20),
```

```
   department varchar(20))
```

```
create type Teacher
```

```
  under Person
```

```
  (salary integer,
```

```
   department varchar(20))
```

- Subtypes can redefine methods by using **overriding method** in place of **method** in the method declaration





## Object-Identity and Reference Types

- Define a type *Department* with a field *name* and a field *head* which is a reference to the type *Person*, with table *people* as scope:

```
create type Department (  
    name varchar (20),  
    head ref (Person) scope people)
```

- We can then create a table *departments* as follows

```
create table departments of Department
```

- We can omit the declaration **scope** *people* from the type declaration and instead make an addition to the **create table** statement:

```
create table departments of Department  
    (head with options scope people)
```



## Path Expressions

- Find the names and addresses of the heads of all departments:

```
select head -> name, head -> address  
from departments
```

- An expression such as “*head* -> *name*” is called a **path expression**
- Path expressions help avoid explicit joins
  - If department head were not a reference, a join of *departments* with *people* would be required to get at the address
  - Makes expressing the query much easier for the user



## Database Models and the Internet

- JDBC
- XML, Xpath, Xquery