
DIGITAL ARITHMETIC

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ABOUT THE BOOK

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Chapter 1: Review of the Basic Number Representations and Algorithms

- General-purpose processors
 - ◇ Main use: numerical computations
 - ◇ Address calculations
- basic operations
- fixed point and floating point
- IEEE standard
- vector processors

- Special-purpose (application-specific) processors
 - for numerically intensive applications
 - single computation or classes of computations

Application-specific processor

- Areas of application:
 - signal processing
 - embedded systems
 - matrix computations
 - graphics, vision, multi-media
 - cryptography and security
 - robotics, instrumentation; others?
- Features:
 - better use of technology
 - improvement in speed, area, power
 - flexibility in
 - * implementation; decomposition into modules
 - * number systems and data formats; algorithms
- Need good design tools; difficult to change; FPGAs?

GENERAL-PURPOSE VS. APPLICATION-SPECIFIC

- Flexibility
- Matching specific applications
- Use of VLSI and special technology
- Use of hardware-level parallel processing
- Lower software overhead

ARITHMETIC PROCESSORS: USER'S VIEW

AP = (operands, operation, results, conditions, singularities)

- Numerical operands and results specified by
 - Set of numerical values $x \in N$ (finite and ordered set)
 - Range $V_{min} \leq x \leq V_{max}$
 - Precision
 - Number representation system (NRS)
- Set of operations: addition, subtraction, multiplication, division
- Conditions: values of the results – zero, negative, etc.
- Singularities: Illegal results – overflow, underflow, Nan, etc.

LEVELS OF DESCRIPTION AND IMPLEMENTATION

- Numerical computations (applications)
- Algorithms
- Arithmetic operations

Operations	Operands-Result	Algorithm
Numerical function	Numbers	Numerical Algorithm
↓	↓	↓
	<i>Number System</i>	<i>Arithmetic design</i>
↓	↓	↓
Digit-vector function	Digit vector	Digit-vector Algorithm
↓	↓	↓
	<i>Digit Coding</i>	<i>RTL design</i>
↓	↓	↓
Bit-vector function	Bit vector	Bit-vector Algorithm
		↓
		<i>Logic Design</i>

NUMBER REPRESENTATION SYSTEMS

value: $x \in N \implies \text{NRS} \implies X \in V$: digit vector

- Digit Vector

$$X = (x_a, \dots, x_i, \dots, x_b)$$

– indexing:

Leftward Zero Origin (LZ) (integers)

$$X = (x_{n-1}, x_{n-2}, \dots, x_0)$$

Rightward One Origin (RO) (fractions)

$$X = (x_1, x_2, \dots, x_n)$$

– Digit set D_i - set of values for digit x_i (usually consecutive)

NUMBER REPRESENTATION (cont.)

- Number of (unambiguously) representable numbers

$$|N| \leq \prod |D_i|$$

- Number representation system

$$F : N \rightarrow V$$

- Choose NRS to
 - allow efficient computation(s)
 - suitable interface with other systems
- Different implementation-performance constraints
 \implies variety of NRS

SOME CHARACTERISTICS OF NRS

a) Range: finite set of digit-vector values

b) Unambiguity: two numbers should not have same representation

$$\text{If } x \in N, y \in N, x \neq y \text{ then } F(x) \neq F(y)$$

c) Nonredundant/redundant

$$\text{Redundant: } F^{-1}(X) = F^{-1}(Y)$$

WEIGHTED NUMBER REPRESENTATION SYSTEM

Integer x represented by digit vector $X = (x_{n-1} \dots, x_0)$,

$$x = \sum_{i=0}^{n-1} x_i \cdot w_i$$

where

$W = (w_{n-1}, \dots, w_0)$ *weight vector*

Define

$R = (r_{n-1}, \dots, r_0)$ *radix vector*

so that

$$w_0 = 1 \quad w_i = w_{i-1} r_{i-1}$$

WEIGHTED NUMBER REPRESENTATION (cont.)

- Fixed-radix NRS

$$r_i = r$$

Then $w_i = r^i$ so that

$$x = \sum_{i=0}^{n-1} x_i r^i$$

- Canonical digit set

$$D_i = \{0, 1, 2, \dots, |r_i| - 1\}$$

- Conventional number system

- Fixed radix positive
- Canonical digit set

NON-CONVENTIONAL FIXED-RADIX SYSTEM

- Negative radix

$$r = -2, x = \sum_{i=0}^{n-1} x_i(-2)^i$$

$$1011 = (-8) + 0 + (-2) + 1 = -9$$

$$0111 = 0 + 4 + (-2) + 1 = 3$$

- Complex radix

$$r = 2j, j = \sqrt{-1}, x_i \in \{0, 1, 2, 3\} \text{ (Knuth's quarter imaginary)}$$

$$W: -8j -4 + 2j + 1$$

$$1231 \implies 1 \times (-8j) + 2 \times (-4) + 3 \times (2j) + 1 \times 1 = -7 - 2j$$

- Non-canonical digit set, $r = 2, \{-1, 0, 1\}$ or $\{0, 1, 2\}$

$$\text{Example: radix 4 } D = \{-3, -2, -1, 0, 1, 2, 3\}$$

$$x = 27 \text{ represented by } (1, 2, 3) \text{ or } (2, -2, 3)$$

REDUNDANT NRS

- Fixed radix r
- Non-canonical digit set

$$D = \{-a, -a + 1, \dots, -1, 0, 1, \dots, b - 1, b\}$$

- Symmetric if $a = b$
- Redundant $a + b + 1 > r$ ($a, b \leq r - 1$)
 - "standard" $a, b \leq r - 1$
 - over-redundant $a, b > r - 1$
 - Redundancy factor

$$\rho = \frac{a}{r - 1}, \quad \rho > \frac{1}{2}$$

EXAMPLES OF REDUNDANT DIGIT SETS

r	a	Digit set	ρ	Comment
2	1	$\{-1, 0, 1\}$	1	minimally/maximally redundant
4	2	$\{-2, -1, 0, 1, 2\}$	$2/3$	minimally redundant
4	3	$\{-3, -2, \dots, 2, 3\}$	1	maximally redundant
4	4	$\{-4, \dots, 4\}$	$4/3$	over-redundant
9	4	$\{-4, \dots, 4\}$	$1/2$	non-redundant
10	5	$\{-5, \dots, 5\}$	$5/9$	minimally redundant
10	6	$\{-6, \dots, 6\}$	$2/3$	redundant
10	9	$\{-9, \dots, 9\}$	1	maximally redundant
10	13	$\{-13, \dots, 13\}$	$13/9$	over-redundant

MIXED-RADIX NUMBER SYSTEM

- $r_i \neq r_j$
- Example: Representation of time $R = (31, 24, 60, 60)$
- Example: Factorial number system

$$\begin{aligned}r_i &= i + 2, \quad i = 0, \dots, n - 1 \\R &= (n + 1, n, \dots, 3, 2) \\w_i &= (i + 1)!\end{aligned}$$

Canonical digit set

Integers in range $0 \leq x \leq (n + 1)! - 1$

NON-WEIGHTED NUMBER SYSTEMS: RESIDUE (RM)

- Base vector B of moduli m_i

$$B = (m_{n-1}, m_{n-2}, \dots, m_0)$$

m_i positive integers and pairwise relatively prime

- Integer x is represented by vector

$$X = (x_{n-1}, x_{n-2}, \dots, x_0)$$

where $x_i = x \bmod m_i$

- Represents uniquely integers in the range

$$0 \leq x < \prod_{i=0}^{n-1} m_i$$

(more later)

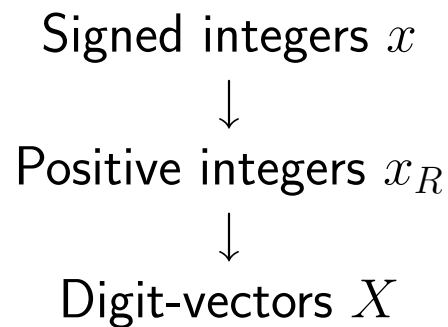
REPRESENTATION OF SIGNED INTEGERS

A. Directly in the number representation

Examples: Signed-Digit Number System

B. With an extra symbol: Sign-and-magnitude

C. Additional mapping on positive integers



Examples:

- True-and-Complement (TC):
 - 2's complement
 - 1s' complement
- Biased representation

TRUE-AND-COMPLEMENT SYSTEM

- $-k \leq x \leq k$ signed integer (implicit value)
- x_R positive integer (representation value)
- C – complementation constant
- Mapping $x_R = x \bmod C$
- Unambiguous if $k < C/2$
- Equivalent to

$$x_R = \begin{cases} x & \mathbf{if} \ x \geq 0 \\ C - |x| & \mathbf{if} \ x < 0 \end{cases}$$

- Converse mapping

$$x = \begin{cases} x_R & \mathbf{if} \ x_R < C/2 \\ x_R - C & \mathbf{if} \ x_R > C/2 \end{cases}$$

NUMBER REPRESENTATION (cont.)

- x_R represented in any number system
- In fixed-radix system two common choices:
 - 2's complement: $C = r^n$ (Range-complement system)
 - 1s' complement: $C = r^n - 1$ (Diminished-radix-complement)

BIASED REPRESENTATION

- B – bias
- $-k \leq x \leq k$
- $x_R = x + B$
- $B \geq k$

TYPES OF ARITHMETIC ALGORITHMS

- Bottom-up development

 - Primitives

 - + Addition/subtraction
 - + Multioperand addition
 - + Arithmetic shifts
 - + Multiplication by digit
 - + Result-digit selection (PLA)
 - + Table look-up
 - + Multiplication

 - Algorithms

 - + Composition of primitives

TYPES (cont.)

- (Digit) Recurrences (continued sums)

- ◇ Residual recurrence: $R[i + 1] = f(R[i], X, Y, Z[i], z_{i+1})$

- Uses: Add/sub, single-position shifts, multiplication by digit

- ◇ Output digit selection: $z_{i+1} = g(R[i], X, Y, Z[i])$

- (keep $R[i + 1]$ bounded)

- Uses: Comparisons, PLA

- ◇ Result recurrence

- $$Z[i + 1] = Z[i] + z_{i+1}r^{i+1} \text{ (continued sum)}$$

- Uses: Concatenation

- Examples:

- ◇ multiplication $R[i + 1] = \frac{1}{r}(R[i] + X \cdot r^n y_i)$

- ◇ division $R[i + 1] = rR[i] - q_{i+1}Y \quad q_{i+1} = g(R[i + 1], Y)$

Types of algorithms

- Continued product recurrences

$$R[i + 1] = f(R[i], X, Y, Z[i], z_{z+1})$$

Uses: Add/sub, variable shifts, mult. by digit

$$z_{i+1} = g(R[i], X, Y, Z[i]) \text{ (keep } R[i + 1] \text{ bounded)}$$

Uses: Comparisons, PLA

$$Z[i + 1] = Z[i](1 + z_{i+1}r^{-(i+1)}) \text{ (continued product)}$$

Uses: Variable shift, addition

- Example:

◇ division

$$R[i + 1] = rR[i](1 + q_{i+1}r^{-(i+1)}) + q_{i+1}$$

$$q_{i+1} = g(R[i], Y)$$

$$Q[i + 1] = Q[i](1 + q_{i+1}r^{-(i+1)})$$

- Iterative Approximations

$$Z[i + 1] = f(Z[i], X, Y) \quad \text{until } g(Z(i)) < \varepsilon$$

Example:

◇ reciprocal

$$Z[i + 1] = Z[i](2 - Z[i]X)$$

- Polynomial Approximations

$$z = a_0 + a_1x + a_2x^2 + \dots$$

PERFORMANCE

- Measures

- + Execution time
- + Throughput

- Improving speed

- a) Arithmetic level

- + Reducing number of steps

 - Example: higher radix

 - Example: combinational instead of sequential

- + Reducing time of step

 - Example: carry-save adder instead of carry-propagate

- + Overlap steps (concurrency/pipelining)

 - Example: multiple generation and addition (in mult.)

 - Example: simultaneous additions (in mult.)

- b) Implementation level

- + Reduce number of logic levels

POWER AND COST

- Measures
 - + Packaging
 - + Interconnection complexity
 - + Number of pins
 - + Number of chips and types of chips
 - + Number of gates and types of gates
 - + Area
 - + Design cost; verification and testing cost
 - + Power dissipation
 - + Power consumption
- Reduction of cost