
Lecture 18

Protection

Goals of Protection

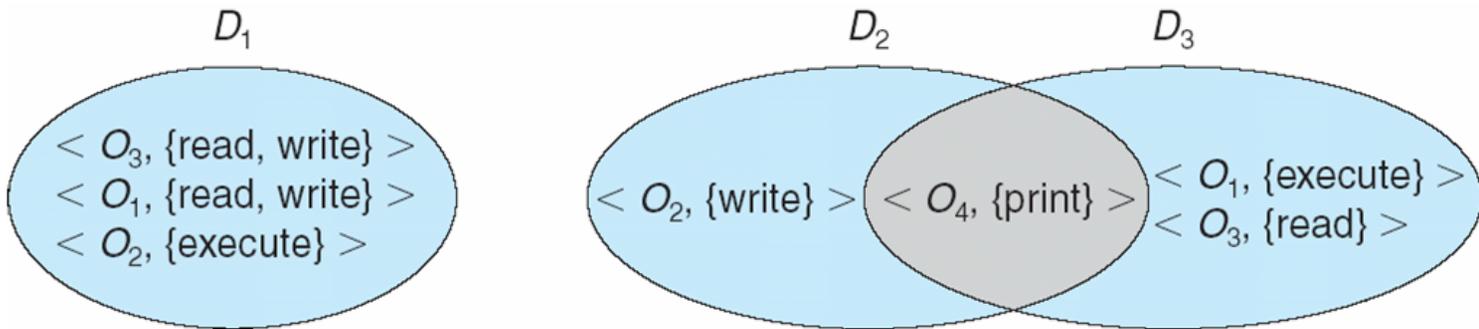
- Protection problem - ensure that each object is accessed correctly and only by those processes that are allowed to do so

Principle of Protection

- Guiding principle – principle of least privilege
 - Programs, users and systems should be given just enough privileges to perform their tasks

Domain Structure

- Access-right = $\langle \text{object-name}, \text{rights-set} \rangle$
where *rights-set* is a subset of all valid operations that can be performed on the object.
- Domain = set of access-rights



Domain Implementation (UNIX)

- System consists of 2 domains:
 - User
 - Supervisor

- UNIX
 - Domain = user-id
 - Domain switch accomplished via file system
 - ▶ Each file has associated with it a domain bit (setuid bit)
 - ▶ When file is executed and setuid = on, then user-id is set to owner of the file being executed. When execution completes user-id is reset

Access Matrix

object domain	F_1	F_2	F_3	printer
D_1	read		read	
D_2				print
D_3		read	execute	
D_4	read write		read write	

Access Matrix

- **Access matrix** design separates mechanism from policy
 - Mechanism
 - ▶ Operating system provides access-matrix + rules
 - ▶ If ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced
 - Policy
 - ▶ User dictates policy
 - ▶ Who can access what object and in what mode

Access Matrix With Domains as Objects

domain \ object	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch
D_3		read	execute					
D_4	read write		read write		switch			

Access Matrix with Copy Rights

object \ domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute		

(a)

object \ domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute	read	

(b)

Access Matrix With *Owner* Rights

domain \ object	F_1	F_2	F_3
D_1	owner execute		write
D_2		read* owner	read* owner write
D_3	execute		

(a)

domain \ object	F_1	F_2	F_3
D_1	owner execute		write
D_2		owner read* write*	read* owner write
D_3		write	write

(b)

Modified Access Matrix

object domain	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch control
D_3		read	execute					
D_4	write		write		switch			

Implementation of Access Matrix

- Each column = Access-control list for one object
Defines who can perform what operation.

Domain 1 = Read, Write
Domain 2 = Read
Domain 3 = Read



- Each Row = Capability List (like a key)
For each domain, what operations allowed on what objects.
 - Object 1 – Read
 - Object 4 – Read, Write, Execute
 - Object 5 – Read, Write, Delete, Copy

Revocation of Access Rights

- **Access List** – Delete access rights from access list
 - Simple, Immediate
- **Capability List** – Scheme required to locate capability in the system before capability can be revoked

Security

The Security Problem

- Security must consider external environment of the system, and protect the system resources
- **Intruders** attempt to breach security
- **Threat** is potential security violation
- **Attack** is attempt to breach security
- Attack can be accidental or malicious, but easier to protect against accidental than malicious misuse

Security Violations

- Categories
 - **Breach of confidentiality**
 - **Breach of integrity**
 - **Breach of availability**
 - **Theft of service**
 - **Denial of service**
- Methods
 - **Masquerading (breach authentication)**
 - **Replay attack**
 - ▶ **Message modification**
 - **Man-in-the-middle attack**
 - **Session hijacking**

Security Measure Levels

- Security must occur at four levels to be effective:
 - **Physical**
 - **Human**
 - ▶ Avoid [social engineering, phishing, dumpster diving](#)
 - **Operating System**
 - **Network**
- Security is as weak as the weakest link in the chain

Program Threats

■ Trojan Horse

- Code segment that misuses its environment
- Exploits mechanisms for allowing programs written by users to be executed by other users
- [Spyware, pop-up browser windows, covert channels](#)

■ Trap Door

- Specific user identifier or password that circumvents normal security procedures
- Could be included in a compiler

■ Logic Bomb

- Program that initiates a security incident under certain circumstances

■ Stack and Buffer Overflow

- Exploits a bug in a program (overflow either the stack or memory buffers)

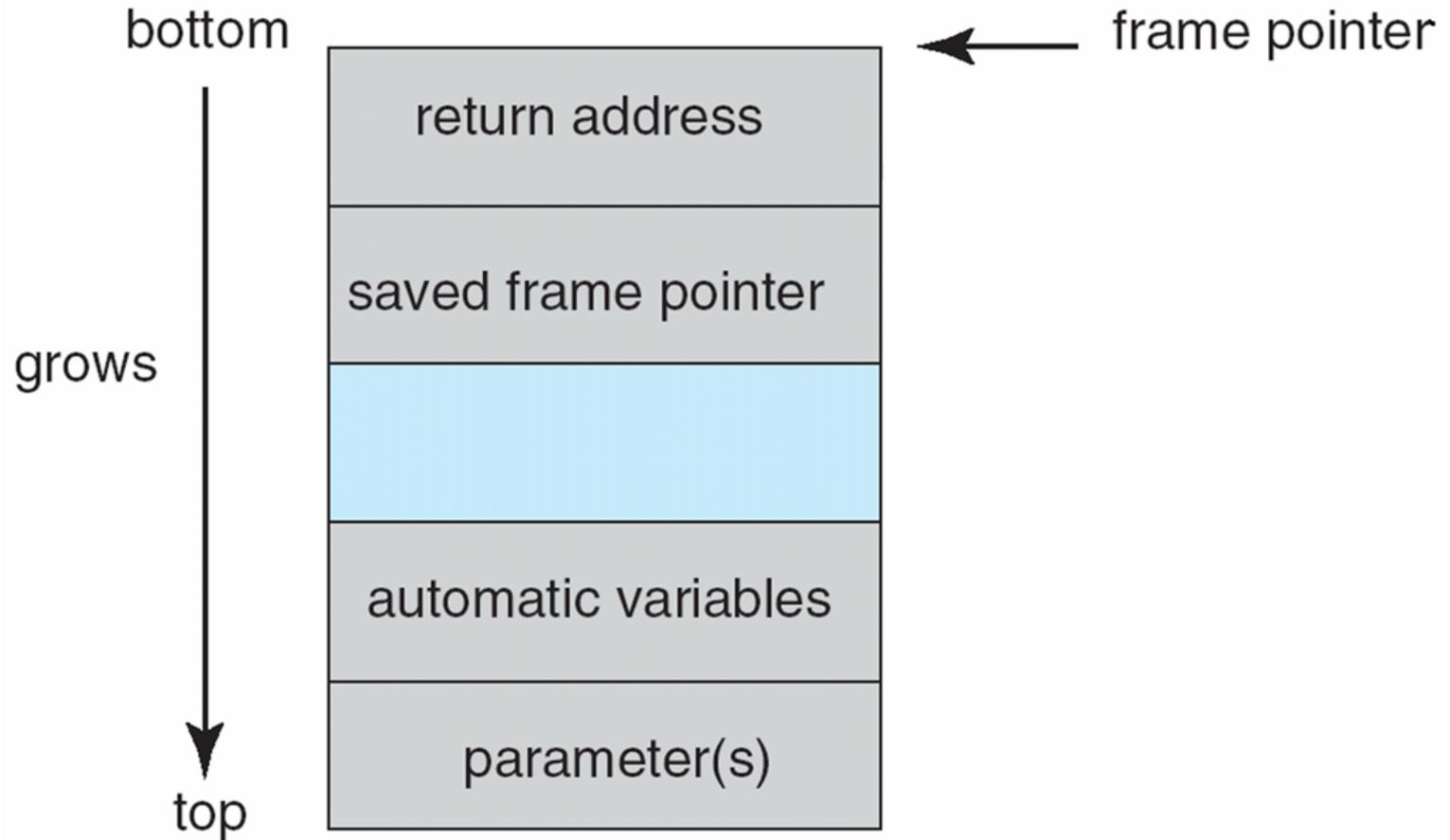
C Program with Buffer-overflow Condition

```
#include <stdio.h>
#define BUFFER SIZE 256
int main(int argc, char *argv[])
{
    char buffer[BUFFER SIZE];
    if (argc < 2)
        return -1;
    else {
        strcpy(buffer, argv[1]);
        return 0;
    }
}
```

See you next time

We will talk about how to play poker online

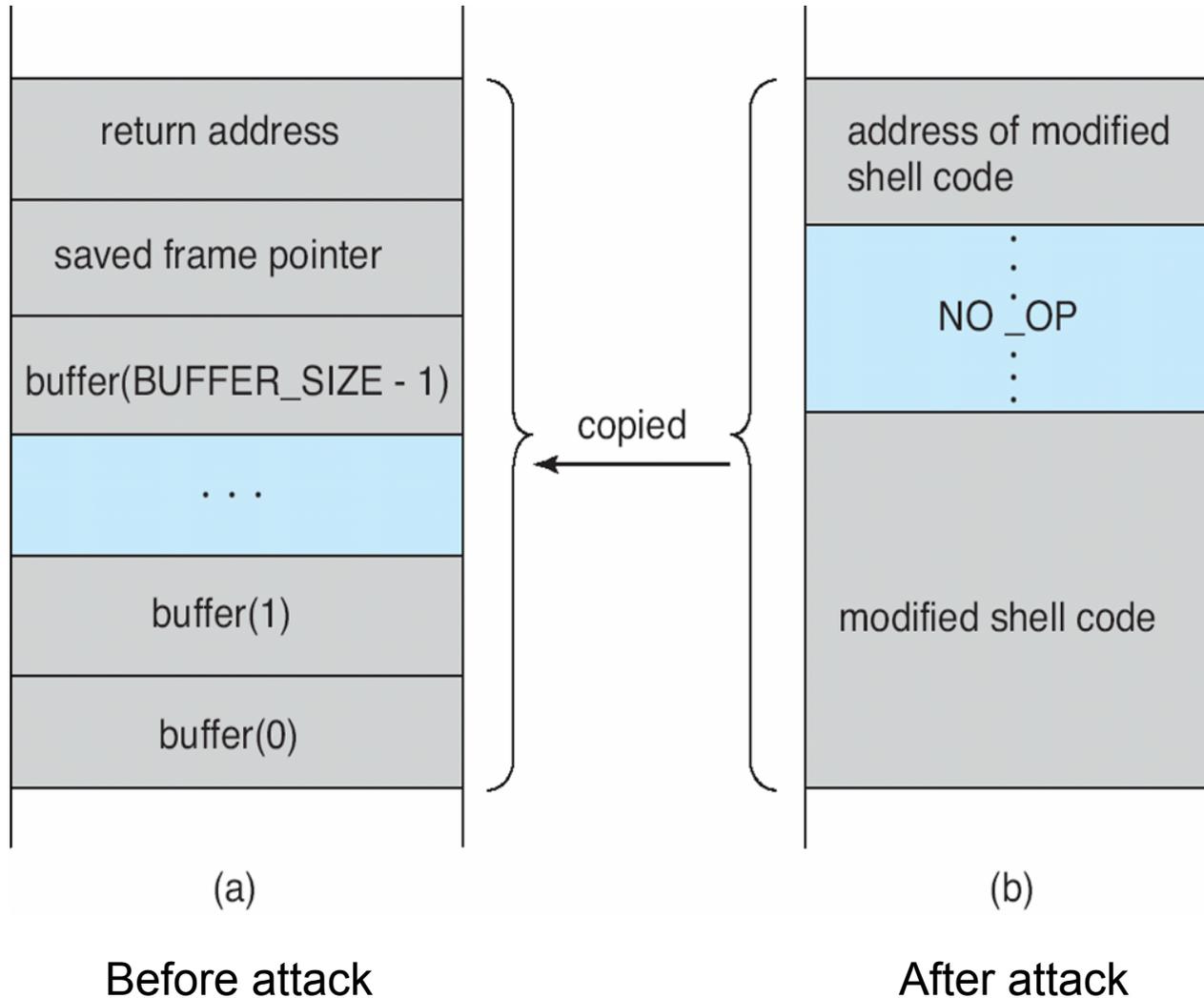
Layout of Typical Stack Frame



Modified Shell Code

```
#include <stdio.h>
int main(int argc, char *argv[])
{
    execvp(``\bin\sh'', ``\bin \sh'', NULL);
    return 0;
}
```

Hypothetical Stack Frame



Program Threats (Cont.)

■ Viruses

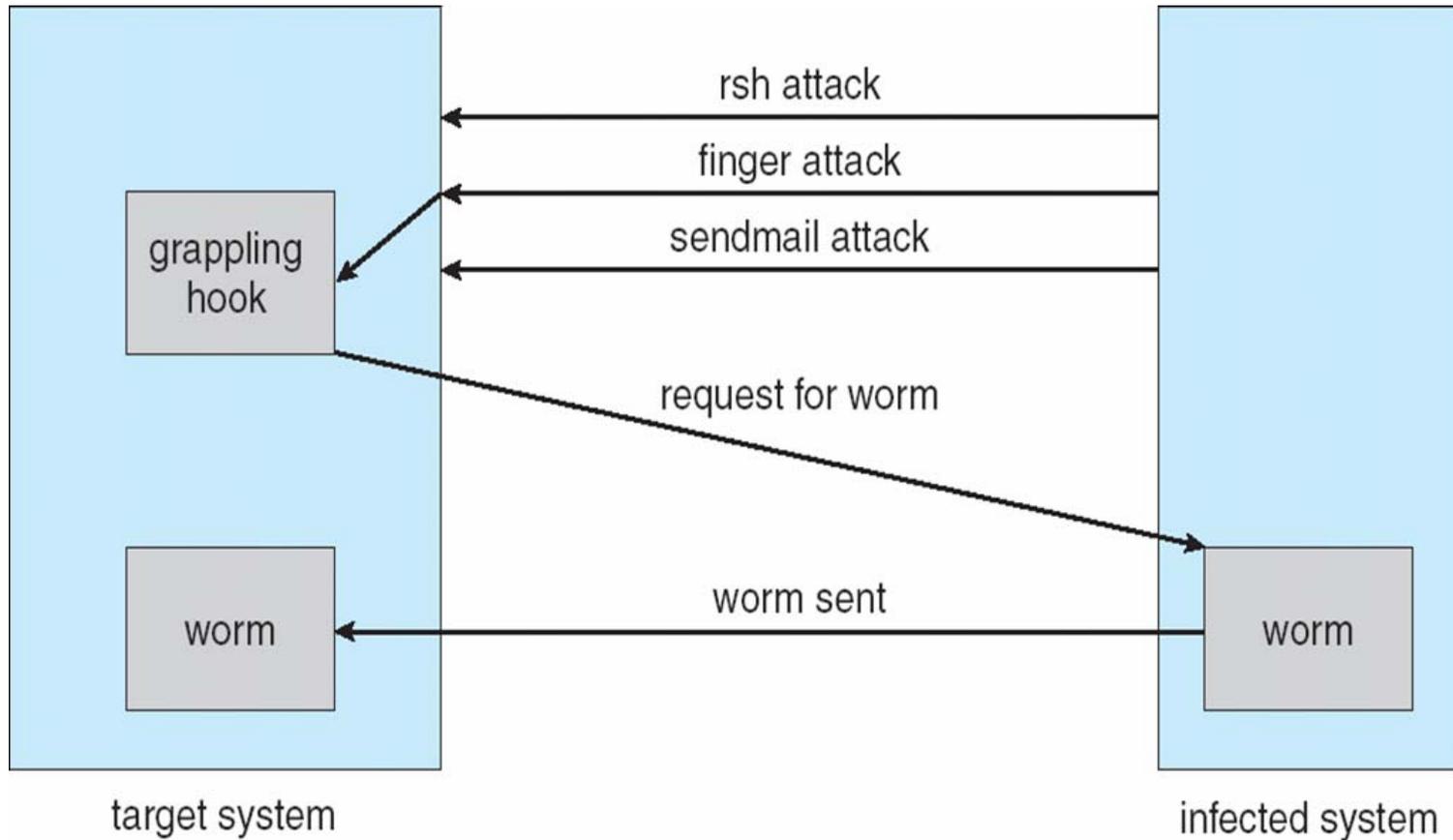
- Code fragment embedded in legitimate program
- Very specific to CPU architecture, operating system, applications
- Usually borne via email or as a macro
 - ▶ Visual Basic Macro to reformat hard drive

```
Sub AutoOpen()  
    Dim oFS  
    Set oFS = CreateObject(''Scripting.FileSystemObject'')  
    vs = Shell(''c:command.com /k format      c:''',vbHide)  
End Sub
```

Program Threats (Cont.)

- **Virus dropper** inserts virus onto the system
- Many categories of viruses, literally many thousands of viruses
 - File
 - Boot
 - Macro
 - Source code
 - Polymorphic
 - Encrypted
 - Stealth
 - Tunneling
 - Multipartite
 - Armored

The Morris Internet Worm



Cryptography as a Security Tool

- Broadest security tool available
 - Source and destination of messages cannot be trusted without cryptography
 - Means to constrain potential senders (*sources*) and / or receivers (*destinations*) of *messages*
- Based on secrets (*keys*)

Encryption

- **Encryption** algorithm consists of
 - Set of K keys
 - Set of M Messages
 - Set of C ciphertexts (encrypted messages)
 - A function $E : K \rightarrow (M \rightarrow C)$. That is, for each $k \in K$, $E(k)$ is a function for generating ciphertexts from messages
 - ▶ Both E and $E(k)$ for any k should be efficiently computable functions
 - A function $D : K \rightarrow (C \rightarrow M)$. That is, for each $k \in K$, $D(k)$ is a function for generating messages from ciphertexts
 - ▶ Both D and $D(k)$ for any k should be efficiently computable functions
- An encryption algorithm must provide this essential property: Given a ciphertext $c \in C$, a computer can compute m such that $E(k)(m) = c$ only if it possesses $D(k)$.
 - Thus, a computer holding $D(k)$ can decrypt ciphertexts to the plaintexts used to produce them, but a computer not holding $D(k)$ cannot decrypt ciphertexts
 - Since ciphertexts are generally exposed (for example, sent on the network), it is important that it be infeasible to derive $D(k)$ from the ciphertexts