## Efficiently and Precisely Locating Memory Leaks and Bloat

Hound: C/C++ Leak & Bloat Detector

## Leak & Bloat ?

#### • Leak

- Reachability leaks(GCable, unreachable)
- Staleness leaks(unused, reachable)
- Bloat
  - Unnecessary excess memory consumption

### Problem: Memory Inefficiency

 Jan/Feb 2008, over 150 leak-related bugs were reported



 Ideal world: Collect high-precision leak reports from real programs with lowoverhead

## Hound: as a solution

- No false positive
- Data Sampling
  - Context-sensitive memory allocation
  - Age-segregated memory allocation
- Virtual compaction





## Why Data Sampling? cont'd





#### High runtime overhead (100X)

## Why Data Sampling? cont'd





False positives

# How to Data Sampling?

• A novel memory manager

• Segregate along 2-D



allocation sites

age

# Site Segregation

void \* houndmalloc (size\_t size) {
// compute hash of calling context.

int context = getContextHash();

Metadata \* m = getMetadata(context);

```
// one more object allocated. m->liveCount++;
```

// use the age-segregated heap to
// satisfy the request, if possible.

if (m->getAgeHeap() != NULL) {

```
return m->getAgeHeap()->malloc (size); }
```

else if (m->getLiveCount() >= 64) {

```
// make a new heap.
m->initAgeHeap();
return m->getAgeHeap()->malloc (size); }
```

#### else {

```
// still below threshold:
// get memory from standard allocator.
return phkmalloc_with_header (size, context);
```

To reduce memory overhead ? Add an extra header word ? Initiate a new age-heap for a site only exceeding a threshold, currently 64 ? Otherwise, allocate objects to a conventional heap

# Age Segregation

```
if (!h->activePage || h->activePage-
>bump == h->activePage->endOfPage) {
  void * page = getNewPage();
  PageEntry * e = createPageEntry (page);
    e->bump = page;
    e->endOfPage = page + PAGE_SIZE;
    e->inUse = 0;
    e->heap = h;
    h->activePage = e;
  }
```



Each heap organizes as a collection of pages.(for each size class) Each page is an array of fixed-sized object slots.

Meta data for each page(bump pointer, # of live objects, a bitmap tracks slots with live objects)

# Age Segregation cont'd

- Keeps all filled pages on aging queue and protects pages on the queue
- Due to cost, cannot protect all pages



## Age Segregation cont'd

- The size of inactive list is controlled adaptively
  - Low runtime overhead & Maximize Useful Info
- Re-evaluate size every I/8 CPU time
  - Page faults > 1.5% of total CPU time  $\rightarrow$  Dec
  - Page faults < 0.5% of total CPU time  $\rightarrow$  Inc
- Low runtime overhead >> useful info

Increase:  $P_I = P_I + max(min(P_A, P_I)/32, 8)$ Decrease:  $P_I = P_I - max(min(P_A, P_I)/8, 8)$ 

## Virtual Compaction

- Why?
  - High fragmentation(potential)
  - Recycles memory from age-segregated heaps only when pages become empty
- To Whom?
  - Toward same sized pages
  - Only for pages have less than 50% occupancy

## Virtual Compaction cont'd

- How?
  - Performing a bitwise AND of several candidate pages' object bitmap
  - Merge them onto a single physical page (mremap call)



 Remap target(physical) page to both virtual pages

## Hound Runtime Overhead

**Hound Runtime Overhead** 



## Hound Memory Overhead

SPECint2006



#### Staleness Computation Accuracy

- Recall(measure the quality of classifier)
  - true positives / (true positives + false negatives)
  - e.g.
    - Consider a report identifies *I* allocation site as the source stale data
    - If this report failed to identify 9 other sites that had stale data
    - recall = 0.1

#### Staleness Computation Accuracy cont'd

**Only underestimate!** 





No false positives

## Questions for Discussion

 Can we appropriately combine two results from Hound and SWAT to make a more decent result?